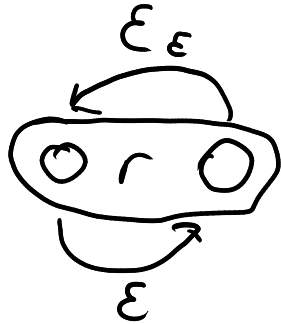
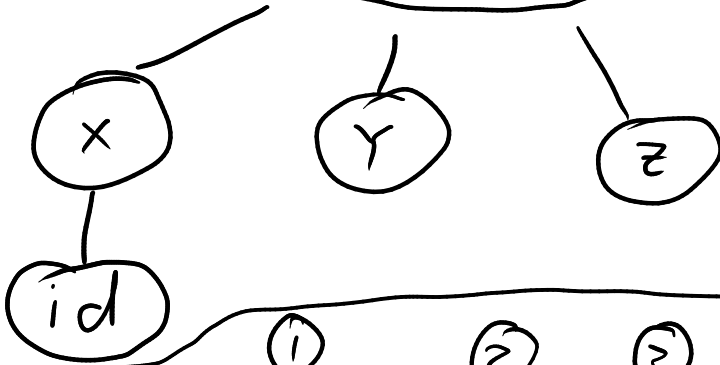


$r^*$



Start nonterminal

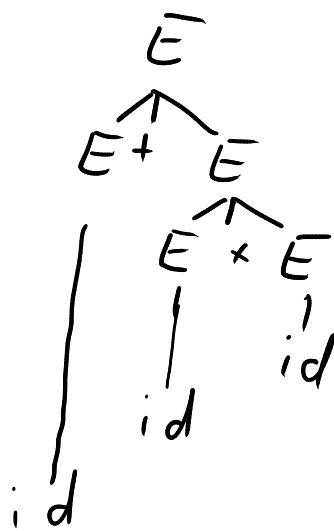


start  $\rightarrow$  XYZ  
 $X \rightarrow$  id

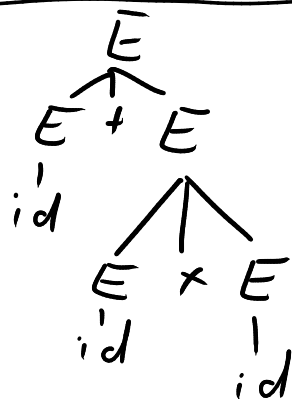
①  $E \rightarrow E + E$   
 rightmost  $E$

id + id x id

- ①  $E + E$
- ②  $E + E \times E$
- ③  $E + E \times id$
- ③  $E + id \times id$
- ③  $id + id \times id$

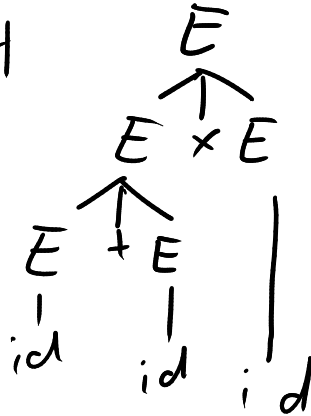


- leftmost  $E$
- ①  $E + E$
  - ③  $id + E$
  - ②  $id + E \times E$
  - ③  $id + id \times E$
  - ③  $id + id \times id$



$\bar{E}$  another leftmost

- ②  $\bar{E} \times E$
- ①  $E + \bar{E} \times E$
- ③  $id + \bar{E} \times \bar{E}$
- ③  $id + id \times E$
- ③  $id + id \times id$



		terminals						X	Y	Z
		a	b	c	d	e	\$			
states	0	shift 5							4	
	1	reduce by rule 7								
	2									
	3									
	4									

```

// LR(1) parsing algorithm
// M is a table like above
// X is the top of the state stack
// a is the next input token
loop {
  if M[X, a] == "shift state j" {
    token_value = advance();
    push_state(j);
    push_value(token_value);
  } else if M[X, a] == "reduce production P" {
    n = new node; // corresponding to the LHS of P
    i = len(rhs(P)) - 1;
    for symbol in reverse(rhs(P)) {
      pop_state();
      n[i] = pop_value();
      i--;
    }
    output(P);
    // X is different from the beginning of the loop; we modified the stack
    push_state(M[X, lhs(P)]);
    push_value(n);
    if P is the augmented start rule {
      break;
    }
  } else {
    // error
  }
}
  
```

} shift

} reduce

```
// LR(0) parsing algorithm
loop {
  if X contains a finished item {
    reduce();
  } else {
    a = advance();
    if M[x, a] == "shift state j" {
      shift(j);
    } else {
      // error
    }
  }
}
}
```