

# LR(1) & LALR(1)

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## LR(k)

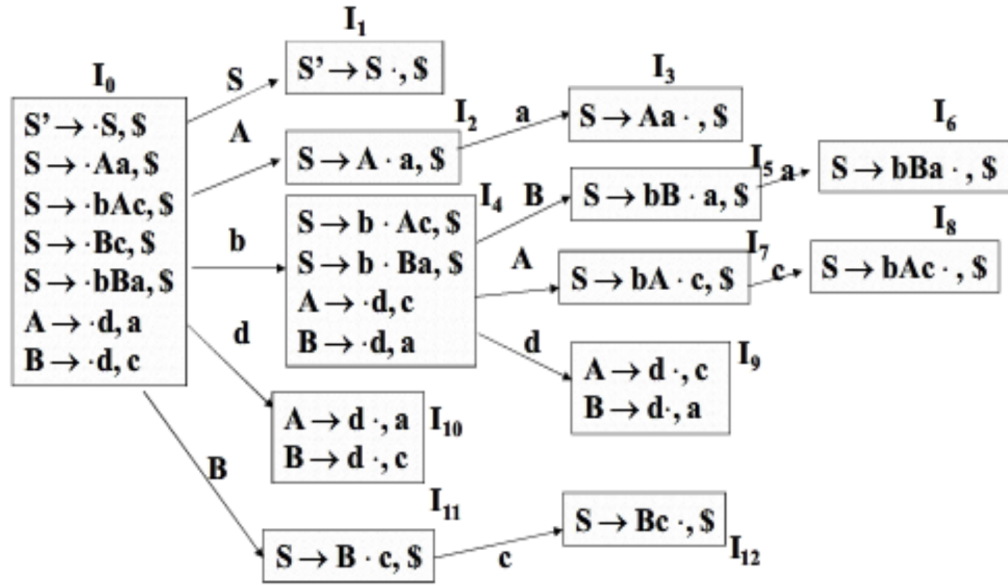
1. Consider the following grammar

$$S \rightarrow A \ a \mid b \ A \ c \mid B \ c \mid b \ B \ a$$

$$A \rightarrow d$$

$$B \rightarrow d$$

- a) Complete LR(1) states for the grammar



b) Complete LR(1) parsing table for the grammar

	<b>a</b>	<b>b</b>	<b>c</b>	<b>d</b>	<b>\$</b>	<b>S</b>	<b>A</b>	<b>B</b>
<b>0</b>		s4		s10		1	2	11
<b>1</b>					acc			
<b>2</b>	s3							
<b>3</b>					r1			
<b>4</b>				s9			7	5
<b>5</b>	S6							
<b>6</b>					r4			
<b>7</b>			s8					
<b>8</b>					r2			
<b>9</b>	r6		r5					
<b>10</b>	r5		r6					
<b>11</b>			s12					
<b>12</b>					r3			

**c) Prove the grammar is not LALR(1)**

Solution: Proof by contradiction. Suppose this is LALR(1) grammar, then the  $state_9$  and  $state_{10}$  can be merged to a single state

$A \rightarrow d, a/c$

$B \rightarrow d, a/c.$

This causes the reduce conflict. Thus, it cannot be a LALR(1) grammar.

d) Which rule can be substituted and removed to make this grammar LALR(1)?

In fact, either  $A \rightarrow d$  or  $B \rightarrow d$  can be substituted and removed. Below is the parsing states for substituting  $B \rightarrow d$ . Clearly, there is no states to be merged. Thus, it is a LALR(1) grammar.

