

The Sensitivity Conjecture

Kunal Chawla

November 16, 2021

Let $f : \{\mathbb{Z}/2\mathbb{Z}\}^n \rightarrow \{0, 1\}$ be a boolean function.

Let $f : \{\mathbb{Z}/2\mathbb{Z}\}^n \rightarrow \{0, 1\}$ be a boolean function.

Definition

The local sensitivity of f at x is

$$s(f, x) := |\{i : f(x + e_i) \neq f(x)\}|.$$

Let $f : \{\mathbb{Z}/2\mathbb{Z}\}^n \rightarrow \{0, 1\}$ be a boolean function.

Definition

The local sensitivity of f at x is

$$s(f, x) := |\{i : f(x + e_i) \neq f(x)\}|.$$

Definition

The sensitivity of f is

$$\max_{x \in \{\mathbb{Z}/2\mathbb{Z}\}^n} s(f, x)$$

Definition

The local block sensitivity, $bs(f, x)$ of f at x is the maximum number of disjoint subsets $B_1, \dots, B_k \subset \{1, \dots, n\}$ such that

$$f(x + \sum_{i \in B_j} e_i) \neq f(x).$$

Definition

The local block sensitivity, $bs(f, x)$ of f at x is the maximum number of disjoint subsets $B_1, \dots, B_k \subset \{1, \dots, n\}$ such that

$$f(x + \sum_{i \in B_j} e_i) \neq f(x).$$

Definition

The block sensitivity, $bs(f)$ is

$$\max_{x \in \{\mathbb{Z}/2\mathbb{Z}\}^n} bs(f, x).$$

The block sensitivity is polynomially related to other complexity measures, such as

The block sensitivity is polynomially related to other complexity measures, such as

- The deterministic decision-tree complexity.

The block sensitivity is polynomially related to other complexity measures, such as

- The deterministic decision-tree complexity.
- The degree of f as a polynomial.

The block sensitivity is polynomially related to other complexity measures, such as

- The deterministic decision-tree complexity.
- The degree of f as a polynomial.
- The certificate complexity of f .

The block sensitivity is polynomially related to other complexity measures, such as

- The deterministic decision-tree complexity.
- The degree of f as a polynomial.
- The certificate complexity of f .

Question

How does block sensitivity relate to sensitivity?

Proposition

For any boolean function f , $bs(f) \geq s(f)$.

Proposition

For any boolean function f , $bs(f) \geq s(f)$.

Proof.

Singletons are sets. □

Proposition

For any boolean function f , $bs(f) \geq s(f)$.

Proof.

Singletons are sets. □

Conjecture (Sensitivity Conjecture)

There exists $C > 0$ such that

$$bs(f) \leq s(f)^C$$

for all boolean functions f .

A related problem

Let Q^n be the graph defined as follows:

- The vertices of Q^n are elements of $(\mathbb{Z}/2\mathbb{Z})^n$
- Two elements $x, y \in \{\mathbb{Z}/2\mathbb{Z}\}^n$ are adjacent if there exists some i with $y = x + e_i$.

A related problem

Let Q^n be the graph defined as follows:

- The vertices of Q^n are elements of $(\mathbb{Z}/2\mathbb{Z})^n$
- Two elements $x, y \in \{\mathbb{Z}/2\mathbb{Z}\}^n$ are adjacent if there exists some i with $y = x + e_i$.

Conjecture (Hypercube Sensitivity)

If H is an induced subgraph of Q^n with $|V(H)| \geq 2^{n-1} + 1$ vertices, then $\Delta(H) \geq \sqrt{n}$.

Hypercube Sensitivity implies Sensitivity

- Let $X_i : (\mathbb{Z}/2\mathbb{Z})^n \rightarrow \{-1, 1\}$ be given by

$$X_i(x) = (-1)^{x_i}.$$

Hypercube Sensitivity implies Sensitivity

- Let $X_i : (\mathbb{Z}/2\mathbb{Z})^n \rightarrow \{-1, 1\}$ be given by

$$X_i(x) = (-1)^{x_i}.$$

- The monomials $\prod_{i \in I} X_i$ are the characters of $(\mathbb{Z}/2\mathbb{Z})^n$, so that every f is a polynomial.

Hypercube Sensitivity implies Sensitivity

- Let $X_i : (\mathbb{Z}/2\mathbb{Z})^n \rightarrow \{-1, 1\}$ be given by

$$X_i(x) = (-1)^{x_i}.$$

- The monomials $\prod_{i \in I} X_i$ are the characters of $(\mathbb{Z}/2\mathbb{Z})^n$, so that every f is a polynomial.
- Let $\mu(x) = X_1(x) \dots X_n(x)$.

Hypercube Sensitivity implies Sensitivity

Proposition

For any boolean function $f : (\mathbb{Z}/2\mathbb{Z})^n \rightarrow \{-1, 1\}$, we have $s(f) \geq \sqrt{\deg(f)}$.

Hypercube Sensitivity implies Sensitivity

Proposition

For any boolean function $f : (\mathbb{Z}/2\mathbb{Z})^n \rightarrow \{-1, 1\}$, we have $s(f) \geq \sqrt{\deg(f)}$.

Proof.

- Let $m = \deg(f)$. By permuting indices, and restricting to $(\mathbb{Z}/2\mathbb{Z})^m$, assume WLOG that $m = n$.

Hypercube Sensitivity implies Sensitivity

Proposition

For any boolean function $f : (\mathbb{Z}/2\mathbb{Z})^n \rightarrow \{-1, 1\}$, we have $s(f) \geq \sqrt{\deg(f)}$.

Proof.

- Let $m = \deg(f)$. By permuting indices, and restricting to $(\mathbb{Z}/2\mathbb{Z})^m$, assume WLOG that $m = n$.
- The Fourier coefficient $X_1 \dots X_n$ is

$$\frac{1}{2^n} \sum_{x \in (\mathbb{Z}/2\mathbb{Z})^n} f(x) \mu(x) \neq 0.$$

Hypercube Sensitivity implies Sensitivity

Proposition

For any boolean function $f : (\mathbb{Z}/2\mathbb{Z})^n \rightarrow \{-1, 1\}$, we have $s(f) \geq \sqrt{\deg(f)}$.

Proof.

- Let $m = \deg(f)$. By permuting indices, and restricting to $(\mathbb{Z}/2\mathbb{Z})^m$, assume WLOG that $m = n$.
- The Fourier coefficient $X_1 \dots X_n$ is

$$\frac{1}{2^n} \sum_{x \in (\mathbb{Z}/2\mathbb{Z})^n} f(x) \mu(x) \neq 0.$$

- Hence $|\{x : \mu(x)f(x) = 1\}| \neq |\{x : \mu(x)f(x) = -1\}|$.

Hypercube Sensitivity implies Sensitivity

Proposition

For any boolean function $f : (\mathbb{Z}/2\mathbb{Z})^n \rightarrow \{-1, 1\}$, we have $s(f) \geq \sqrt{\deg(f)}$.

Proof.

- Let $m = \deg(f)$. By permuting indices, and restricting to $(\mathbb{Z}/2\mathbb{Z})^m$, assume WLOG that $m = n$.
- The Fourier coefficient $X_1 \dots X_n$ is

$$\frac{1}{2^n} \sum_{x \in (\mathbb{Z}/2\mathbb{Z})^n} f(x) \mu(x) \neq 0.$$

- Hence $|\{x : \mu(x)f(x) = 1\}| \neq |\{x : \mu(x)f(x) = -1\}|$. Let E be the larger set.

Hypercube Sensitivity implies Sensitivity

Proposition

For any boolean function $f : (\mathbb{Z}/2\mathbb{Z})^n \rightarrow \{-1, 1\}$, we have $s(f) \geq \sqrt{\deg(f)}$.

Proof.

- Let $m = \deg(f)$. By permuting indices, and restricting to $(\mathbb{Z}/2\mathbb{Z})^m$, assume WLOG that $m = n$.
- The Fourier coefficient $X_1 \dots X_n$ is

$$\frac{1}{2^n} \sum_{x \in (\mathbb{Z}/2\mathbb{Z})^n} f(x) \mu(x) \neq 0.$$

- Hence $|\{x : \mu(x)f(x) = 1\}| \neq |\{x : \mu(x)f(x) = -1\}|$. Let E be the larger set.



Hypercube Sensitivity implies Sensitivity

Proposition

For any boolean function $f : (\mathbb{Z}/2\mathbb{Z})^n \rightarrow \{-1, 1\}$, we have $s(f) \geq \sqrt{\deg(f)}$.

Proof.

- By Hypercube Sensitivity, there exists $x \in E$ such that $x + e_i \in E$ for $\lceil \sqrt{n} \rceil$ values of i .

Hypercube Sensitivity implies Sensitivity

Proposition

For any boolean function $f : (\mathbb{Z}/2\mathbb{Z})^n \rightarrow \{-1, 1\}$, we have $s(f) \geq \sqrt{\deg(f)}$.

Proof.

- By Hypercube Sensitivity, there exists $x \in E$ such that $x + e_i \in E$ for $\lceil \sqrt{n} \rceil$ values of i .
- As $\mu(x + e_i) = -\mu(x)$ we have $f(x + e_i) \neq f(x)$, for at least \sqrt{n} values of i .



Why is \sqrt{n} sharp?

Theorem

There exists an induced subgraph $H \subset Q^n$ such that $|V(H)| = 2^{n-1} + 1$ and $\Delta(H) \leq \lceil \sqrt{n} \rceil$.

Why is \sqrt{n} sharp?

Theorem

There exists an induced subgraph $H \subset Q^n$ such that $|V(H)| = 2^{n-1} + 1$ and $\Delta(H) \leq \lceil \sqrt{n} \rceil$.

Proof.

- Let $\mu(x) := (-1)^{x_1 + \dots + x_n}$.

Why is \sqrt{n} sharp?

Theorem

There exists an induced subgraph $H \subset Q^n$ such that $|V(H)| = 2^{n-1} + 1$ and $\Delta(H) \leq \lceil \sqrt{n} \rceil$.

Proof.

- Let $\mu(x) := (-1)^{x_1 + \dots + x_n}$.
- Let I_1, \dots, I_k be a partition of $\{1, \dots, n\}$.

Why is \sqrt{n} sharp?

Theorem

There exists an induced subgraph $H \subset Q^n$ such that $|V(H)| = 2^{n-1} + 1$ and $\Delta(H) \leq \lceil \sqrt{n} \rceil$.

Proof.

- Let $\mu(x) := (-1)^{x_1 + \dots + x_n}$.
- Let I_1, \dots, I_k be a partition of $\{1, \dots, n\}$. Let

$$V_j = \{x : x_i = 0 \text{ for } i \in I_j\}.$$

Why is \sqrt{n} sharp?

Theorem

There exists an induced subgraph $H \subset Q^n$ such that $|V(H)| = 2^{n-1} + 1$ and $\Delta(H) \leq \lceil \sqrt{n} \rceil$.

Proof.

- Let $\mu(x) := (-1)^{x_1 + \dots + x_n}$.
- Let I_1, \dots, I_k be a partition of $\{1, \dots, n\}$. Let

$$V_j = \{x : x_i = 0 \text{ for } i \in I_j\}.$$

- By the binomial theorem,

$$\sum_{x \in \bigcap_{j \in J} V_j} \mu(x) = 1_{n = \sum_{j \in J} |I_j|}.$$

Proof.

- By inclusion-exclusion,

$$\sum_{x \in (\mathbb{Z}/2\mathbb{Z})^n \setminus \bigcup_{j=1}^k V_j} \mu(x) = (-1)^k.$$

Proof.

- By inclusion-exclusion,

$$\sum_{x \in (\mathbb{Z}/2\mathbb{Z})^n \setminus \bigcup_{j=1}^k V_j} \mu(x) = (-1)^k.$$

- Therefore

$$\sum_{x \in \bigcup_{j=1}^k V_j} \mu(x) = (-1)^{k+1}.$$

Proof.

- By inclusion-exclusion,

$$\sum_{x \in (\mathbb{Z}/2\mathbb{Z})^n \setminus \bigcup_{j=1}^k V_j} \mu(x) = (-1)^k.$$

- Therefore

$$\sum_{x \in \bigcup_{j=1}^k V_j} \mu(x) = (-1)^{k+1}.$$

- So

$$\sum_{x \in \bigcup_{j=1}^k V_j} (-1)^{k+1} \mu(x) + \sum_{x \notin \bigcup_{j=1}^k V_j} (-1)^k \mu(x) = 2.$$



Proof.

- Let

$$H_1 = \{x \notin \cup_{j=1}^k V_j \text{ such that } (-1)^k \mu(x) = 1\}.$$

Proof.

- Let

$$H_1 = \{x \notin \cup_{j=1}^k V_j \text{ such that } (-1)^k \mu(x) = 1\}.$$

- Let

$$H_2 = \{x \in \cup_{j=1}^k V_j \text{ such that } (-1)^{k+1} \mu(x) = 1\}.$$

Proof.

- Let

$$H_1 = \{x \notin \cup_{j=1}^k V_j \text{ such that } (-1)^k \mu(x) = 1\}.$$

- Let

$$H_2 = \{x \in \cup_{j=1}^k V_j \text{ such that } (-1)^{k+1} \mu(x) = 1\}.$$

- Let $H = H_1 \cup H_2$, so that $|E| = 2^{n-1} + 1$.



Proof.

- If $x \in H_2$ then $(-1)^k \mu(x + e_i) = 1$ for all i .

Proof.

- If $x \in H_2$ then $(-1)^k \mu(x + e_i) = 1$ for all i . Also, $x + e_i \in V_j$ as long as $i \in I_j$.

Proof.

- If $x \in H_2$ then $(-1)^k \mu(x + e_i) = 1$ for all i . Also, $x + e_i \in V_j$ as long as $i \in I_j$. Hence $\deg(x) \leq \max_j |I_j|$.

Proof.

- If $x \in H_2$ then $(-1)^k \mu(x + e_i) = 1$ for all i . Also, $x + e_i \in V_j$ as long as $i \in I_j$. Hence $\deg(x) \leq \max_j |I_j|$.
- If $x \in H_1$ then $(-1)^{k+1} \mu(x + e_i) = 1$, and there is at most one i such that $x + e_i \in V_j$.

Proof.

- If $x \in H_2$ then $(-1)^k \mu(x + e_i) = 1$ for all i . Also, $x + e_i \in V_j$ as long as $i \in I_j$. Hence $\deg(x) \leq \max_j |I_j|$.
- If $x \in H_1$ then $(-1)^{k+1} \mu(x + e_i) = 1$, and there is at most one i such that $x + e_i \in V_j$. Hence $\deg(x) \leq k$.

Proof.

- If $x \in H_2$ then $(-1)^k \mu(x + e_i) = 1$ for all i . Also, $x + e_i \in V_j$ as long as $i \in I_j$. Hence $\deg(x) \leq \max_j |I_j|$.
- If $x \in H_1$ then $(-1)^{k+1} \mu(x + e_i) = 1$, and there is at most one i such that $x + e_i \in V_j$. Hence $\deg(x) \leq k$.
- This implies $\Delta(H) \leq \max\{\max_j |I_j|, k\}$.

Proof.

- If $x \in H_2$ then $(-1)^k \mu(x + e_i) = 1$ for all i . Also, $x + e_i \in V_j$ as long as $i \in I_j$. Hence $\deg(x) \leq \max_j |I_j|$.
- If $x \in H_1$ then $(-1)^{k+1} \mu(x + e_i) = 1$, and there is at most one i such that $x + e_i \in V_j$. Hence $\deg(x) \leq k$.
- This implies $\Delta(H) \leq \max\{\max_j |I_j|, k\}$.
- Let I_j be a partition of $[n]$ into $\sim \sqrt{n}$ pieces of size $\sim \sqrt{n}$, so that $\Delta(H) \leq \lceil \sqrt{n} \rceil$.



We've seen

- Why Hypercube Sensitivity implies the Sensitivity conjecture.

We've seen

- Why Hypercube Sensitivity implies the Sensitivity conjecture.
- Why the \sqrt{n} is sharp

Theorem

If H is an induced subgraph of Q^n with $|V(H)| \geq 2^{n-1} + 1$ vertices, then $\Delta(H) \geq \sqrt{n}$.

Lemma

Let H be a graph. If λ is an eigenvalue of the adjacency matrix $A(H)$, then $|\lambda| \leq \Delta(H)$.

Lemma

Let H be a graph. If λ is an eigenvalue of the adjacency matrix $A(H)$, then $|\lambda| \leq \Delta(H)$.

Proof.

Let v be an eigenvector for λ and suppose that v_1 is the largest coordinate of v (in absolute value).

Lemma

Let H be a graph. If λ is an eigenvalue of the adjacency matrix $A(H)$, then $|\lambda| \leq \Delta(H)$.

Proof.

Let v be an eigenvector for λ and suppose that v_1 is the largest coordinate of v (in absolute value). Then

$$|\lambda v_1| = \left| \sum_{k=1}^m A_{1,k} v_k \right| \leq \sum_{k=1}^m |A_{1,k}| |v_k| \leq \Delta(H) |v_1|.$$



Lemma

Let H be a graph. If λ is an eigenvalue of the adjacency matrix $A(H)$, then $|\lambda| \leq \Delta(H)$.

Proof.

Let v be an eigenvector for λ and suppose that v_1 is the largest coordinate of v (in absolute value). Then

$$|\lambda v_1| = \left| \sum_{k=1}^m A_{1,k} v_k \right| \leq \sum_{k=1}^m |A_{1,k}| |v_k| \leq \Delta(H) |v_1|.$$



We have actually proven the following:

Lemma

Let B be a matrix such that $|B| \leq A(H)$ entrywise. Then if λ is an eigenvalue of B , we have $|\lambda| \leq \Delta(H)$.

The Strategy Now

Construct an auxiliary object with nice spectral properties.

The Strategy Now

Construct an auxiliary object with nice spectral properties.

Question

How can we estimate the eigenvalues of a pseudo-adjacency matrix of H ?

Theorem (Cauchy's interlace Theorem)

Suppose that A is a symmetric $n \times n$ matrix, and B is a principal $m \times m$ submatrix, where $m < n$. If the eigenvalues of A are $\lambda_1 \leq \dots \leq \lambda_n$ and the eigenvalues of B are $\beta_1 \leq \dots \leq \beta_m$, then

$$\lambda_i \leq \beta_i \leq \lambda_{i+n-m}.$$

Lemma (Min-Max principle)

Suppose A is an $n \times n$ symmetric matrix, and let $R_A(x) = \frac{\langle Ax, x \rangle}{\langle x, x \rangle}$ be the Rayleigh quotient. If $\lambda_1 \leq \dots \leq \lambda_k \leq \dots \leq \lambda_n$ are the eigenvalues of A , then

$$\lambda_k = \min_U \{ \max_x \{ R_A(x) \mid x \in U \text{ and } x \neq 0 \mid \dim(U) = k \} \},$$

and

$$\lambda_k = \max_U \{ \min_x \{ R_A(x) \mid x \in U \text{ and } x \neq 0 \mid \dim(U) = n - k + 1 \} \}.$$

Lemma (Min-Max principle)

Suppose A is an $n \times n$ symmetric matrix, and let $R_A(x) = \frac{\langle Ax, x \rangle}{\langle x, x \rangle}$ be the Rayleigh quotient. If $\lambda_1 \leq \dots \leq \lambda_k \leq \dots \leq \lambda_n$ are the eigenvalues of A , then

$$\lambda_k = \min_U \{ \max_x \{ R_A(x) \mid x \in U \text{ and } x \neq 0 \mid \dim(U) = k \} \},$$

and

$$\lambda_k = \max_U \{ \min_x \{ R_A(x) \mid x \in U \text{ and } x \neq 0 \mid \dim(U) = n - k + 1 \} \}.$$

Proof.

Count dimensions □

Theorem (Cauchy's interlace Theorem)

Suppose that A is a symmetric $n \times n$ matrix, and B is a principal $m \times m$ submatrix, where $m < n$. If the eigenvalues of A are $\lambda_1 \leq \dots \leq \lambda_n$ and the eigenvalues of B are $\beta_1 \leq \dots \leq \beta_m$, then

$$\lambda_i \leq \beta_i \leq \lambda_{i+n-m}.$$

Proof.

Let $A = \begin{bmatrix} B & X^T \\ X & Z \end{bmatrix}$, and let $\{v_k, \dots, v_n\}$ be the last $n - k + 1$ eigenvectors of A , and $\{w_1, \dots, w_m\}$ the eigenvectors of B .

Proof.

Let $A = \begin{bmatrix} B & X^T \\ X & Z \end{bmatrix}$, and let $\{v_k, \dots, v_n\}$ be the last $n - k + 1$ eigenvectors of A , and $\{w_1, \dots, w_m\}$ the eigenvectors of B .

Let $V = \text{span}\{v_k, \dots, v_n\}$, $W = \text{span}(w_1, \dots, w_k)$. Lift $\widehat{W} = \begin{pmatrix} W \\ 0 \end{pmatrix}$.

Proof.

Let $A = \begin{bmatrix} B & X^T \\ X & Z \end{bmatrix}$, and let $\{v_k, \dots, v_n\}$ be the last $n - k + 1$ eigenvectors of A , and $\{w_1, \dots, w_m\}$ the eigenvectors of B .

Let $V = \text{span}\{v_k, \dots, v_n\}$, $W = \text{span}(w_1, \dots, w_k)$. Lift $\widehat{W} = \begin{pmatrix} W \\ 0 \end{pmatrix}$.

Since $\dim V = n - k + 1$ and $\dim \widehat{W} = k$, there exists some $\widehat{w} \in V \cap \widehat{W}$.
Hence $R_A(\widehat{w}) = R_B(w)$

Proof.

Let $A = \begin{bmatrix} B & X^T \\ X & Z \end{bmatrix}$, and let $\{v_k, \dots, v_n\}$ be the last $n - k + 1$ eigenvectors of A , and $\{w_1, \dots, w_m\}$ the eigenvectors of B .

Let $V = \text{span}\{v_k, \dots, v_n\}$, $W = \text{span}(w_1, \dots, w_k)$. Lift $\widehat{W} = \begin{pmatrix} W \\ 0 \end{pmatrix}$.

Since $\dim V = n - k + 1$ and $\dim \widehat{W} = k$, there exists some $\widehat{w} \in V \cap \widehat{W}$.
Hence $R_A(\widehat{w}) = R_B(w)$

Hence

$$\lambda_k \leq \min_{x \in V} R_A(x) \leq R_A(\widehat{w}) = R_B(w) \leq \max_{w \in W} R_B(x) = \beta_k.$$

Proof.

Let $A = \begin{bmatrix} B & X^T \\ X & Z \end{bmatrix}$, and let $\{v_k, \dots, v_n\}$ be the last $n - k + 1$ eigenvectors of A , and $\{w_1, \dots, w_m\}$ the eigenvectors of B .

Let $V = \text{span}\{v_k, \dots, v_n\}$, $W = \text{span}(w_1, \dots, w_k)$. Lift $\widehat{W} = \begin{pmatrix} W \\ 0 \end{pmatrix}$.

Since $\dim V = n - k + 1$ and $\dim \widehat{W} = k$, there exists some $\widehat{w} \in V \cap \widehat{W}$.
Hence $R_A(\widehat{w}) = R_B(w)$

Hence

$$\lambda_k \leq \min_{x \in V} R_A(x) \leq R_A(\widehat{w}) = R_B(w) \leq \max_{w \in W} R_B(x) = \beta_k.$$

The other direction is similar, with $V = \text{span}(v_1, \dots, v_{k+n-m})$ and $W = \text{span}(w_k, \dots, w_m)$. □

Proof of the main theorem

Consider the following sequence of matrices given by

$$A_1 = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \text{ and } A_n = \begin{bmatrix} A_{n-1} & I \\ I & -A_{n-1} \end{bmatrix}.$$

Proof of the main theorem

Consider the following sequence of matrices given by

$$A_1 = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \text{ and } A_n = \begin{bmatrix} A_{n-1} & I \\ I & -A_{n-1} \end{bmatrix}.$$

Lemma

Replacing every -1 with 1 , we get the adjacency matrix of Q^n .

Proof of the main theorem

Consider the following sequence of matrices given by

$$A_1 = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \text{ and } A_n = \begin{bmatrix} A_{n-1} & I \\ I & -A_{n-1} \end{bmatrix}.$$

Lemma

$$A_n^2 = nI$$

Proof of the main theorem

Consider the following sequence of matrices given by

$$A_1 = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} \text{ and } A_n = \begin{bmatrix} A_{n-1} & I \\ I & -A_{n-1} \end{bmatrix}.$$

Lemma

The eigenvalues of A_n are $\pm\sqrt{n}$ with multiplicity 2^{n-1} .

Theorem

If H is an induced subgraph of Q^n with $|V(H)| \geq 2^{n-1} + 1$, then $\Delta H \geq \sqrt{n}$.

Theorem

If H is an induced subgraph of Q^n with $|V(H)| \geq 2^{n-1} + 1$, then $\Delta H \geq \sqrt{n}$.

Proof.

Let A_H be the induced principle submatrix of A .

Theorem

If H is an induced subgraph of Q^n with $|V(H)| \geq 2^{n-1} + 1$, then $\Delta H \geq \sqrt{n}$.

Proof.

Let A_H be the induced principle submatrix of A . As $|H| \geq 2^{n-1} + 1$, then by Cauchy's interlacing theorem, we have

$$\Delta H \geq \lambda_1(A_H) \geq \lambda_{2^{n-1}}(A_n) = \sqrt{n}.$$



In Summary

We

- 1 Translated a problem into a combinatorial statement about graphs.
- 2 Encountered a variational representation of eigenvalues.
- 3 Used spectral methods to lower bound the maximum degree of a graph.